Distributed Optimization of Traffic Signals in Arterial Networks

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Background - Traffic Congestion

- Negative Impact
 - Traffic delay
 - Fuel consumption
 - Air pollution
- Approaches
 - Congestion pricing
 - Road expansion
 - Traffic signal control





Small = less than 500,000 Medium = 500,000 to 1 million Large = 1 million to 3 million Very Large = more than 3 million

2019 URBAN MOBILITY REPORT, published by The Texas A&M Transportation Institute.



Traffic congestion in America

https://www.usnews.com/news/slideshows/ worst-traffic-cities-in-america-ranked

Deterministic Model - Overview

- Spatial decomposition
 - Partition road into cells.
 - Types of cells: origin (*O*), ordinary (*E*),
 diverge (*V*), intersection (*I*), merge (*M*),
 destination (*D*).
- Temporal decomposition
 - Divide time horizon in time steps as $\{1, \dots, T\}$.



* Daganzo, C. (1992). The cell transmission model. part II: Network traffic.



Deterministic Model – Variables, Parameters, Objective Function and Constraints

- Relax dynamic equation constraints
 - Transform them into linear constraints
 - Add term in objective function to maximize the flow of vehicles.
- Final objective function
 - Maximize throughput of network
 - Maximize the flow of vehicles

min
$$-\sum_{c \in \mathcal{D}} \sum_{t=1}^{T} n_{ct} - \alpha \sum_{c \in \mathcal{C}} \sum_{t=1}^{T} (T-t) y_{ct}$$

- Variables
- *y*: Number of vehicles leaving cell
- *n*: Number of vehicles inside cell
- *z*: Indicator of which cycle and phase of green time this time step is in
- Parameters
- Q: Flow capacity
- β : Turning ratio
- *W*: Ratio between shock-wave propagation speed and the flow-free speed
- *N*: Jam density
- D: Demand
- *n*^{*init*}: Initialized number of vehicles
- *N_{cy}*: Number of cycles
- $\boldsymbol{\mathcal{R}}$: Set of intersections

Deterministic Model – Flow-balance Constraints

• Flow-balance constraints are different for different types of cells.



Deterministic Model – Signal Constraints

- Optimize green time, cycle length and offset.
- Phase sequence: $1 \rightarrow 2 \rightarrow 3 \rightarrow 4$.



- Phase 1 Phase 2 Phase 3
- Use integer variables to describe 'if-then' constraints.
- Variables and parameters *z*: Indicator integer variable *b*, *e*: Begin and end time of green phase *g*: Green time, *l*: Cycle length, *o*: Offset U, ϵ : Sufficient large and small parameters for 'if-then' constraints *G_{min}*, *G_{max}*: Minimum and maximum green time

 $-U \cdot z_{1ijj't} + \epsilon \leq t - e_{ijj'} \leq U(1 - z_{1ijj't}), \ \forall i \in \mathcal{R}, \ j' = 1, \cdots, N_{cn}, \ t = 1, \cdots, T$ $-U \cdot z_{2ijj't} + \epsilon \leq b_{ijj'} - t \leq U(1 - z_{2ijj't}), \forall i \in \mathcal{R}, \forall j \in \mathcal{F}, t = 1, \cdots, T$ $\sum z_{1ijj't} + z_{2ijj't} \le 5, \ j' = 1, \cdots, N_{cy}, \ t = 1, \cdots, T$

 $o_i \leq l_i, \ \forall i \in \mathcal{R}$

 $b_{i1j'} = l_{i1j'} \cdot j' - o_i \ \forall i \in \mathcal{R}, \ j' = 1, \cdots, N_{cy}$ $e_{i1j'} = b_{i1j'} + g_{i1}, \ \forall i \in \mathcal{R}, \ j' = 1, \cdots, N_{cy}$ $b_{i2j'} = e_{i1j'}, \forall i \in \mathcal{R}, j' = 1, \cdots, N_{cy}$ $e_{i2j'} = b_{i2j'} + g_{i2}, \ \forall i \in \mathcal{R}, \ j' = 1, \cdots, N_{cu}$ $b_{i3j'} = e_{i2j'}, \forall i \in \mathcal{R}, j' = 1, \cdots, N_{cy}$ $e_{i3i'} = b_{i3i'} + g_{i3}, \ \forall i \in \mathcal{R}, \ j' = 1, \cdots, N_{cu}$ $b_{i4j'} = e_{i3j'}, \forall i \in \mathcal{R}, j' = 1, \cdots, N_{cu}$ $e_{i4i'} = b_{i4i'} + g_{i4}, \ \forall i \in \mathcal{R}, \ j' = 1, \cdots, N_{cu}$ $l_i = \sum g_{ij}, \ \forall i \in \mathcal{R}$

 $G_{\min} \leq g_{ij} \leq G_{\max}, \ \forall i \in \mathcal{R}, \ \forall j \in \mathcal{F}$

 $z_{1ijj't}, z_{2ijj't} \in \{0, 1\}, \forall i \in \mathcal{R}, \forall j \in \mathcal{F}, j' = 1, \cdots, N_{cu}, t = 1, \cdots, T$

If time step *t* is during green time $[b_{ijj'}, e_{ijj'}]$, then $z_{1ijj't} =$ $z_{2ijj't} = 1.$

Computation of start and end time of green phase given green time, offset and cycle length.

Minimum and maximum green time constraints.

Deterministic Model – Distributed Formulation

- Input boundary cells: receiving inflow from cell of neighboring area.
- Output boundary cells: sending outflow to cell of neighboring area.



Algorithm – Alternating Direction Method of Multipliers (ADMM)

• Lagrangian function

$$\begin{aligned} \mathcal{L}(\mathbf{x}, \mathbf{u}, \mathbf{v}, \boldsymbol{\kappa}, \boldsymbol{\lambda}, \boldsymbol{\mu}, \boldsymbol{\nu}) &= \sum_{i=1}^{N} \left(-\sum_{c \in \mathcal{D}_{i}} \sum_{t=1}^{T} n_{ct} - \alpha \sum_{c \in \mathcal{C}_{i}} \sum_{t=1}^{T} (T - t) y_{ct} \right) \\ &+ \frac{L}{2} \sum_{i=1}^{N_{I}} (\|\mathbf{B}_{i} \mathbf{x}_{i} + \mathbf{v}_{i} - \mathbf{b}_{i} + \kappa_{i}\|_{2}^{2} + \|\mathbf{x}_{i} - \mathbf{u}_{i} + \boldsymbol{\nu}_{i}\|_{2}^{2} \\ &+ \sum_{c \in \mathcal{B}_{i}^{O}} \sum_{t=1}^{T_{o}} \left(y_{ct} + s_{ct} + W u_{n_{dct}} - W N_{ct} + \lambda_{ct} \right)^{2} \\ &+ \sum_{c \in \mathcal{B}_{i}^{I}} \sum_{t=1}^{T_{o}} \left(n_{ct+1} - n_{ct} + y_{ct} - u_{y_{pct}} + \mu_{ct} \right)^{2} \right) \end{aligned}$$

- Algorithm steps
 - Step 1: Minimize Lagrangian function with respect to *x*.

 $\mathbf{x}^{l+1} = \mathrm{argmin}_{\mathbf{x} \in \mathcal{C}_x} \mathcal{L}(\mathbf{x}, \mathbf{u}^l, \mathbf{v}^l, \boldsymbol{\kappa}^l, \boldsymbol{\lambda}^l, \boldsymbol{\mu}^l, \boldsymbol{\nu}^l)$

• Step 2: Minimize Lagrangian function with respect to *z*.

$$\mathbf{z}^{l+1} = \mathrm{argmin}_{\mathbf{z} \in \mathcal{C}_z} \mathcal{L}(\mathbf{x}^{l+1}, \mathbf{u}, \mathbf{v}, \boldsymbol{\kappa}^l, \boldsymbol{\lambda}^l, \boldsymbol{\mu}^l, \boldsymbol{\nu}^l)$$

• Step 3: Update dual variables.

 $\begin{aligned} \kappa_{i}^{l+1} &= \mathbf{B}_{i} \mathbf{x}_{i}^{l+1} + \mathbf{u}_{i}^{l+1} - \mathbf{b}_{i} + \kappa_{i}^{l} \\ \lambda_{ct}^{l+1} &= y_{ct}^{l+1} + s_{ct}^{l+1} + W u_{n_{dct}}^{l+1} - W N_{ct} + \lambda_{ct}^{l}, \ c \in \mathcal{B}_{i}^{O} \\ \mu_{ct}^{l+1} &= n_{ct+1}^{l+1} - n_{ct}^{l+1} + y_{ct}^{l+1} - u_{y_{pct}}^{l+1} + \mu_{ct}^{l}, \ c \in \mathcal{B}_{i}^{I} \\ \boldsymbol{\nu}_{i}^{l+1} &= \mathbf{x}_{i}^{l+1} - \mathbf{u}_{i}^{l+1} + \boldsymbol{\nu}_{i}^{l} \end{aligned}$

Algorithm – Alternating Direction Method of Multipliers with Heuristic

• Distributed form

$$\begin{array}{ll} \min & \sum_{i=1}^{N} \left(-\sum_{c \in \mathcal{D}_{i}} \sum_{t=1}^{T} n_{ct} - \alpha \sum_{c \in \mathcal{C}_{i}} \sum_{t=1}^{T} (T-t) y_{ct} \right) \\ \text{s.t.} & \mathbf{A}_{i} \mathbf{x}_{i} = \mathbf{a}_{i}, \ \forall i \in \mathcal{R} \\ & \mathbf{B}_{i} \mathbf{x}_{i} + \mathbf{v}_{i} = \mathbf{b}_{i}, \ \forall i \in \mathcal{R} \\ & y_{ct} + s_{ct} = WN_{ct} - Wu_{n_{d_{c}t}}, \ \forall i \in \mathcal{R}, \ \forall c \in \mathcal{B}_{i}^{O}, \ t = 1, \cdots, T \\ & n_{ct+1} = n_{ct} + u_{y_{pct}} - y_{ct}, \ \forall i \in \mathcal{R}, \ \forall c \in \mathcal{B}_{i}^{I}, \ t = 1, \cdots, T \\ & \mathbf{u}_{i}^{(1)} = \mathbf{x}_{i}, \ \mathbf{u}_{i}^{(2)} = \mathbf{x}_{i}, \ \forall i \in \mathcal{R} \\ & \mathbf{u}_{z}^{(1)} \in [0, 1]^{8N_{I}N_{cy}T}, \ \mathbf{u}_{z}^{(2)} \in \{u : \|u - \frac{1}{2}\mathbf{1}_{8N_{I}N_{cy}T}\|_{2}^{2} = 2N_{I}N_{cy}T\} \\ & \mathbf{v}_{i} \geq 0, \ \forall i \in \mathcal{R} \\ & \mathbf{x}_{i}^{lb} \leq \mathbf{x}_{i} \leq \mathbf{x}^{ub}, \ \forall i \in \mathcal{R} \\ & \mathbf{z}_{1ijj't}, \ z_{2ijj't} \in \{0, 1\}, \ \forall i \in \mathcal{R}, \ \forall j \in \mathcal{F}, \ j' = 1, \cdots, N_{cy}, \ t = 1, \cdots, T \end{array}$$

• Variable update

• Undate of integer variables are different

Stochastic Model

- Stochastic parameters: turning ratio (β) and demand (D).
- Assumption: known distribution of stochastic parameters.
- First-stage variables decisions made "here and now"
 - Begin and end time of green phase (b, e), green time (g), offset (o), cycle length (l) and corresponding integer variables (z_1, z_2) .
- Second-stage variables "wait and see" recourse decisions
 - Number of vehicles leaving (y) and inside each cell (n).
- Two-stage stochastic program formulation
 - Use θ to estimate the objective value of second stage problem.
 - First-stage problem (master problem) min $\sum_{k=1}^{K} p^k \theta^k$
 - s.t. Signal constraints and previous added cuts
 - Second-stage problem $\theta^k = \min -\sum_{c \in D} \sum_{t=1}^T n_{ct}^k \alpha \sum_{c \in C} \sum_{t=1}^T (T-t) y_{ct}^k$ (subproblem)
 - s.t. Flow constraints for each scenario

Algorithm – Distributed Benders Cut with Cycle Estimation

- Benders cut for each iteration, add constraints of θ and first-stage variables based on the dual solution of subproblem to master problem. $\theta^k \ge F_D^k(z_1, z_2, \hat{\rho}^k, \hat{\sigma}^k, \hat{\pi}^k, \hat{\gamma}^k, \hat{\delta}^k, \hat{\tau}^k)$
- Distributed Benders cut
 - Signal constraints are separate for each intersection. Solve master problem for each intersection separately.

$$\min \quad \sum_{k=1}^{K} p^k \theta_i^k$$

- s.t. Signal constraints and previous added cuts for each intersection
- The constraint in Benders cut can be written as summation of variables of each intersection. Add constraints to master problem of each intersection separately. $\theta_i^k \ge F_{D_i}^k(z_{1i}, z_{2i}, \hat{\rho}^k, \hat{\sigma}^k, \hat{\pi}^k, \hat{\gamma}^k, \hat{\delta}^k, \hat{\tau}^k)$
- Estimation of cycle reduce the number of integer variables
 - Use the cycle length of previous iteration to estimate the cycle each time step is in.
 - Before estimation: consider z_1, z_2 for every cycle $j' = 1, \dots, N_{cy}$ to indicate if the time step is in cycle j'.
 - After estimation: consider z_1, z_2 for only two cycles $j' = \lfloor t/l \rfloor, \lfloor t/l \rfloor$ to indicate if the time step is in cycle j'.

Result

- Settings
 - Grid network size: 4×4 , time horizon: 200s, number of scenarios: 10
 - Flow capacity: Q = 2 for intersection cells, Q = 4 for other cells.
 - Jam density: N = 8 for intersection cells, N = 16 for other cells.
 - Minimum green time: 12s, maximum green time: 32s
 - Demand: Poisson distribution with mean randomly generated from [0.4,1.2] (west-east direction) and [0.2, 0.6] (south-north direction)
 - Turning ratio: Randomly selected from given possible turning ratio set
- Problem size
 - 12800 integer variables, 501728 continuous variables, 1176448 constraints.
- Performance
 - Gurobi: no feasible solution with 7200s time limit.
 - Benders cut: cannot solve master problem with 7200s time limit after 1 iteration.
 - Distributed Benders cut: Obtain solution after 40 iterations (14000s).
 - Distributed Benders cut with cycle estimation: Obtain solution after 40 iterations (9500s).



Result of Distributed Benders Cut with Cycle Estimation

• Grid network size: 4 × 4, time horizon: 200s (50 time steps), number of scenarios: 10



- Average time of solving master problem: 35s
- Average time of solving sub problem: 163s
- Average time of each iteration: 198s

Original Signal Control Plan Provided by Distributed Benders Cut with Cycle Estimation

- Each line is corresponding to an intersection.
- Different color segment indicates the time period of different phases.



Signal control plan of intersections

- Each graph is corresponding to a corridor.
- The moving direction of corridor is from west to east.



- Each graph is corresponding to a corridor.
- The moving direction of corridor is from east to west.

- Each graph is corresponding to a corridor.
- The moving direction of corridor is from north to south.

- Each graph is corresponding to a corridor.
- The moving direction of corridor is from south to north.

Future Work

- Apply model and algorithm to solve larger instance.
- Employ parallel computing to solve problem in each iteration.

Thank you! Questions?