# Distributed Optimization of Traffic Signals in Arterial Networks

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# Outline

- Background
	- Traffic Congestion
- Deterministic Model
	- Overview
	- Variables, Parameters, Objective Function and Flow-balance constraints
	- Signal Constraints
- Algorithm for Deterministic Model
	- Distributed Form of Deterministic Model
	- Alternating Direction Method of Multipliers Algorithm
- Stochastic Model
- Algorithm for Stochastic Model
	- Distributed Benders Cut with cycle estimation
- Numerical Results

# Background - Traffic Congestion

- Negative Impact
	- Traffic delay
	- Fuel consumption
	- Air pollution
- Approaches
	- Congestion pricing
	- Road expansion
	- Traffic signal control 2019 URBAN MOBILITY REPORT,





Small = less than  $500,000$ Medium =  $500,000$  to 1 million

Large =  $1$  million to  $3$  million Very Large = more than 3 million

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#### Traffic congestion in America

[https://www.usnews.com/news/slideshows/](https://www.usnews.com/news/slideshows/worst-traffic-cities-in-america-ranked) [worst-traffic-cities-in-america-ranked](https://www.usnews.com/news/slideshows/worst-traffic-cities-in-america-ranked)

### Deterministic Model - Overview

- Spatial decomposition
	- Partition road into cells.
	- Types of cells: origin  $(0)$ , ordinary  $(\mathcal{E})$ , diverge  $(\mathcal{V})$ , intersection  $(\mathcal{I})$ , merge  $(\mathcal{M})$ , destination  $(D)$ .
- Temporal decomposition
	- Divide time horizon in time steps as  $\{1, \dots, T\}$ .



\* Daganzo, C. (1992). The cell transmission model. part II: Network traffic.



# Deterministic Model – Variables, Parameters, Objective Function and Constraints

- Relax dynamic equation constraints
	- Transform them into linear constraints
	- Add term in objective function to maximize the flow of vehicles.
- Final objective function
	- Maximize throughput of network
	- Maximize the flow of vehicles

$$
\min \quad -\sum_{c \in \mathcal{D}} \sum_{t=1}^{T} n_{ct} - \alpha \sum_{c \in \mathcal{C}} \sum_{t=1}^{T} (T-t) y_{ct}
$$

- Variables
- : Number of vehicles leaving cell
- : Number of vehicles inside cell
- : Indicator of which cycle and phase of green time this time step is in
- Parameters
- : Flow capacity
- $\beta$ : Turning ratio
- W: Ratio between shock-wave propagation speed and the flow-free speed
- N: Jam density
- D: Demand
- $n^{\mathit{init}}$ : Initialized number of vehicles
- $N_{cv}$ : Number of cycles
- : Set of intersections

#### Deterministic Model – Flow-balance Constraints

• Flow-balance constraints are different for different types of cells.



 $y_{ct} > 0$ ,  $n_{ct} > 0$ ,  $\forall c \in \mathcal{C}, t = 1, \cdots, T$ 

# Deterministic Model – Signal Constraints

- Optimize green time, cycle length and offset.
- Phase sequence:  $1 \rightarrow 2 \rightarrow 3 \rightarrow 4$ .



- Phase 1 Phase 2 Phase 3 Phase 4
- Use integer variables to describe 'if-then' constraints.
- Variables and parameters : Indicator integer variable  $b, e$ : Begin and end time of green phase  $g:$  Green time,  $l:$  Cycle length,  $o:$  Offset  $U, \epsilon$ : Sufficient large and small parameters for 'if-then' constraints  $G_{min}$ ,  $G_{max}$ : Minimum and maximum green time

 $-U \cdot z_{1i\,i\,j'\,t} + \epsilon \le t - e_{i\,i\,j'} \le U(1-z_{1i\,i\,j'\,t}), \ \forall i \in \mathcal{R}, \ \ j' = 1, \cdots, N_{cu}, \ t = 1, \cdots, T$  $-U \cdot z_{2ij}t + \epsilon \leq b_{ij}t - t \leq U(1 - z_{2ij}t), \ \forall i \in \mathcal{R}, \ \forall j \in \mathcal{F}, \ t = 1, \cdots, T$  $\sum z_{1ijj't} + z_{2ijj't} \le 5, \ j' = 1, \cdots, N_{cy}, \ t = 1, \cdots, T$ 

 $o_i \leq l_i, \ \forall i \in \mathcal{R}$ 

 $b_{i1j'} = l_{i1j'} \cdot j' - o_i \; \forall i \in \mathcal{R}, \; j' = 1, \cdots, N_{cu}$  $e_{i1j'} = b_{i1j'} + g_{i1}, \ \forall i \in \mathcal{R}, \ j' = 1, \cdots, N_{cu}$  $b_{i2j'} = e_{i1j'}, \ \forall i \in \mathcal{R}, \ j' = 1, \cdots, N_{cu}$  $e_{i2j'} = b_{i2j'} + q_{i2}, \ \forall i \in \mathcal{R}, \ j' = 1, \cdots, N_{cu}$  $b_{i3j'} = e_{i2j'}, \ \forall i \in \mathcal{R}, \ j' = 1, \cdots, N_{cy}$  $e_{i3j'} = b_{i3j'} + g_{i3}, \ \forall i \in \mathcal{R}, \ j' = 1, \cdots, N_{cy}$  $b_{i4j'} = e_{i3j'}$ ,  $\forall i \in \mathcal{R}, j' = 1, \cdots, N_{cu}$  $e_{i4j'} = b_{i4j'} + g_{i4}, \ \forall i \in \mathcal{R}, \ j' = 1, \cdots, N_{cu}$  $l_i = \sum g_{ij}, \ \forall i \in \mathcal{R}$ 

 $G_{\min} \le g_{ij} \le G_{\max}, \ \forall i \in \mathcal{R}, \ \forall j \in \mathcal{F}$ 

 $z_{1iji't}, z_{2iji't} \in \{0,1\}, \forall i \in \mathcal{R}, \forall j \in \mathcal{F}, j' = 1, \cdots, N_{cu}, t = 1, \cdots, T$ 

If time step  $t$  is during green time  $[b_{ijj'}, e_{ijj'}]$ , then  $z_{1ijj't} =$  $z_{2ijj't} = 1.$ 

Computation of start and end time of green phase given green time, offset and cycle length.

Minimum and maximum green time constraints.

#### Deterministic Model – Distributed Formulation

- Input boundary cells: receiving inflow from cell of neighboring area.
- Output boundary cells: sending outflow to cell of neighboring area.



# Algorithm – Alternating Direction Method of Multipliers (ADMM)

• Lagrangian function • Algorithm steps

$$
\mathcal{L}(\mathbf{x}, \mathbf{u}, \mathbf{v}, \kappa, \lambda, \mu, \nu) = \sum_{i=1}^{N} \left( -\sum_{c \in \mathcal{D}_i} \sum_{t=1}^{T} n_{ct} - \alpha \sum_{c \in \mathcal{C}_i} \sum_{t=1}^{T} (T - t) y_{ct} \right) \n+ \frac{L}{2} \sum_{i=1}^{N_I} (||\mathbf{B}_i \mathbf{x}_i + \mathbf{v}_i - \mathbf{b}_i + \kappa_i||_2^2 + ||\mathbf{x}_i - \mathbf{u}_i + \nu_i||_2^2) \n+ \sum_{c \in \mathcal{B}_i^O} \sum_{t=1}^{T_o} (y_{ct} + s_{ct} + W u_{n_{dct}} - W N_{ct} + \lambda_{ct})^2 \n+ \sum_{c \in \mathcal{B}_i^I} \sum_{t=1}^{T_o} (n_{ct+1} - n_{ct} + y_{ct} - u_{y_{pc}t} + \mu_{ct})^2)
$$

- - Step 1: Minimize Lagrangian function with respect to  $x$ .

 $\mathbf{x}^{l+1} = \operatorname{argmin}_{\mathbf{x} \in C_x} \mathcal{L}(\mathbf{x}, \mathbf{u}^l, \mathbf{v}^l, \mathbf{\kappa}^l, \boldsymbol{\lambda}^l, \boldsymbol{\mu}^l, \boldsymbol{\nu}^l)$ 

• Step 2: Minimize Lagrangian function with respect to  $z$ .

$$
\mathbf{z}^{l+1} = \mathrm{argmin}_{\mathbf{z} \in \mathcal{C}_z} \mathcal{L}(\mathbf{x}^{l+1}, \mathbf{u}, \mathbf{v}, \kappa^l, \boldsymbol{\lambda}^l, \boldsymbol{\mu}^l, \boldsymbol{\nu}^l)
$$

• Step 3: Update dual variables.

 $\kappa_i^{l+1} = B_i x_i^{l+1} + u_i^{l+1} - b_i + \kappa_i^l$  $\lambda_{ct}^{l+1} = y_{ct}^{l+1} + s_{ct}^{l+1} + Wu_{n_{det}}^{l+1} - WN_{ct} + \lambda_{ct}^{l}, c \in \mathcal{B}_{i}^{O}$  $\mu_{ct}^{l+1} = n_{ct+1}^{l+1} - n_{ct}^{l+1} + y_{ct}^{l+1} - u_{y_{net}}^{l+1} + \mu_{ct}^{l}, \ c \in \mathcal{B}_{i}^{I}$  $\nu_i^{l+1} = \mathbf{x}_i^{l+1} - \mathbf{u}_i^{l+1} + \nu_i^l$ 

# Algorithm – Alternating Direction Method of Multipliers with Heuristic

• Distributed form • Variable update

$$
\min \sum_{i=1}^{N} \left( -\sum_{c \in \mathcal{D}_i} \sum_{t=1}^{T} n_{ct} - \alpha \sum_{c \in \mathcal{C}_i} \sum_{t=1}^{T} (T-t) y_{ct} \right)
$$
\n
$$
\text{s.t.} \quad \mathbf{A}_i \mathbf{x}_i = \mathbf{a}_i, \ \forall i \in \mathcal{R}
$$
\n
$$
\mathbf{B}_i \mathbf{x}_i + \mathbf{v}_i = \mathbf{b}_i, \ \forall i \in \mathcal{R}
$$
\n
$$
\mathbf{B}_i \mathbf{x}_i + \mathbf{v}_i = \mathbf{b}_i, \ \forall i \in \mathcal{R}
$$
\n
$$
\mathbf{B}_i \mathbf{x}_i + \mathbf{v}_i = \mathbf{b}_i, \ \forall i \in \mathcal{R}
$$
\n
$$
\mathbf{B}_i \mathbf{v}_i + s_{ct} = W N_{ct} - W u_{n_{dc}}, \ \forall i \in \mathcal{R}, \ \forall c \in \mathcal{B}_i^O, \ t = 1, \dots, T
$$
\n
$$
\mathbf{a}_i^{(1)} = \mathbf{x}_i, \ \mathbf{u}_i^{(2)} = \mathbf{x}_i, \ \forall i \in \mathcal{R}
$$
\n
$$
\mathbf{u}_i^{(2)} \in [0, 1]^{8N_I N_{cy} T}, \ \mathbf{u}_i^{(2)} \in \{u : \|u - \frac{1}{2}1_{8N_I N_{cy} T}\|_2^2 = 2N_I N_{cy} T\}
$$
\n
$$
\mathbf{U}_i \geq 0, \ \forall i \in \mathcal{R}
$$
\n
$$
\mathbf{x}^{lb} \leq \mathbf{x}_i \leq \mathbf{x}^{ub}, \ \forall i \in \mathcal{R}
$$
\n
$$
\mathbf{x}^{lb} \leq \mathbf{x}_i \leq \mathbf{x}^{ub}, \ \forall i \in \mathcal{R}
$$
\n
$$
\mathbf{x}^{lb} \leq \mathbf{x}_i \leq \mathbf{x}^{ub}, \ \forall i \in \mathcal{R}
$$
\n
$$
\mathbf{x}^{lb} \leq \mathbf{x}_i \leq \mathbf{x}^{lb}, \ \forall i \in \mathcal{R}
$$
\n
$$
\mathbf{x}^{lb} \leq \math
$$

• Undate of integer variables are different

### Stochastic Model

- Stochastic parameters: turning ratio  $(\beta)$  and demand  $(D)$ .
- Assumption: known distribution of stochastic parameters.
- First-stage variables decisions made "here and now"
	- Begin and end time of green phase  $(b, e)$ , green time  $(g)$ , offset  $(o)$ , cycle length  $(l)$  and corresponding integer variables  $(z_1, z_2)$ .
- Second-stage variables "wait and see" recourse decisions
	- Number of vehicles leaving  $(y)$  and inside each cell  $(n)$ .
- Two-stage stochastic program formulation
	- Use  $\theta$  to estimate the objective value of second stage problem.
	- First-stage problem  $K$ (master problem)
		- min  $\sum p^k \theta^k$
		- s.t. Signal constraints and previous added cuts
	- Second-stage problem  $\theta^k = \min$   $-\sum_{c \in \mathcal{D}} \sum_{t=1}^t n_{ct}^k \alpha \sum_{c \in \mathcal{C}} \sum_{t=1}^t (T-t) y_{ct}^k$ (subproblem)
		- s.t. Flow constraints for each scenario

# Algorithm – Distributed Benders Cut with Cycle Estimation

- Benders cut for each iteration, add constraints of  $\theta$  and first-stage variables based on the dual solution of subproblem to master problem.  $\theta^k \ge F_D^k(z_1, z_2, \hat{\rho}^k, \hat{\sigma}^k, \hat{\tau}^k, \hat{\gamma}^k, \hat{\delta}^k, \hat{\tau}^k)$
- Distributed Benders cut
	- Signal constraints are separate for each intersection. Solve master problem for each intersection separately.

$$
\min \sum_{k=1}^{K} p^k \theta_i^k
$$

- s.t. Signal constraints and previous added cuts for each intersection
- The constraint in Benders cut can be written as summation of variables of each intersection. Add constraints to master problem of each intersection separately.  $\theta_i^k \geq F_{Di}^k(z_{1i}, z_{2i}, \hat{\rho}^k, \hat{\sigma}^k, \hat{\tau}^k, \hat{\gamma}^k, \hat{\delta}^k, \hat{\tau}^k)$
- Estimation of cycle reduce the number of integer variables
	- Use the cycle length of previous iteration to estimate the cycle each time step is in.
	- Before estimation: consider  $z_1, z_2$  for every cycle  $j' = 1, \dots, N_{cy}$  to indicate if the time step is in cycle  $j'$ .
	- After estimation: consider  $z_1, z_2$  for only two cycles  $j' = \lfloor t/l \rfloor$ ,  $\lceil t/l \rceil$  to indicate if the time step is in cycle  $i'$ .

# Result

- Settings
	- Grid network size:  $4 \times 4$ , time horizon: 200s, number of scenarios: 10
	- Flow capacity:  $Q = 2$  for intersection cells,  $Q = 4$  for other cells.
	- Jam density:  $N = 8$  for intersection cells,  $N = 16$  for other cells.
	- Minimum green time: 12s, maximum green time: 32s
	- Demand: Poisson distribution with mean randomly generated from [0.4,1.2] (west-east direction) and [0.2, 0.6 (south -north direction)
	- Turning ratio: Randomly selected from given possible turning ratio set
- Problem size
	- 12800 integer variables, 501728 continuous variables, 1176448 constraints.
- Performance
	- Gurobi: no feasible solution with 7200s time limit.
	- Benders cut: cannot solve master problem with 7200s time limit after 1 iteration.
	- Distributed Benders cut: Obtain solution after 40 iterations (14000s).
	- Distributed Benders cut with cycle estimation: Obtain solution after 40 iterations (9500s).



# Result of Distributed Benders Cut with Cycle Estimation

• Grid network size:  $4 \times 4$ , time horizon: 200s (50 time steps), number of scenarios: 10



- Average time of solving master problem: 35s
- Average time of solving sub problem: 163s
- Average time of each iteration: 198s

# Original Signal Control Plan Provided by Distributed Benders Cut with Cycle Estimation

- Each line is corresponding to an intersection.
- Different color segment indicates the time period of different phases.



Signal control plan of intersections

- Each graph is corresponding to a corridor.
- The moving direction of corridor is from west to east.



- Each graph is corresponding to a corridor.
- The moving direction of corridor is from east to west.



- Each graph is corresponding to a corridor.
- The moving direction of corridor is from north to south.



- Each graph is corresponding to a corridor.
- The moving direction of corridor is from south to north.



#### Future Work

- Apply model and algorithm to solve larger instance.
- Employ parallel computing to solve problem in each iteration.

Thank you! Questions?