Connected Vehicles Based Traffic Signal Timing Optimization

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Wan Li, Xuegang (Jeff) Ban*

Abstract—We study the traffic signal control problem with connected vehicles (CVs) by assuming a fixed cycle length so that the proposed model can be extended readily for the coordination of multiple signals. The problem can be first formulated as a mixed-integer nonlinear program, by considering information of individual vehicle’s trajectories (i.e., second-by-second vehicle locations and speeds) and their realistic driving/car-following behaviors. The objective function is to minimize the weighted sum of total fuel consumption and travel time. Due to the large dimension of the problem and the complexity of the nonlinear car-following model, solving the nonlinear program directly is challenging. We then reformulate the problem as a Dynamic Programming (DP) model by dividing the timing decisions into stages (one stage for a signal phase) and approximating the fuel consumption and travel time of a stage as functions of the state and decision variables of the stage. We also propose a two-step method to make sure that the obtained optimal solution can lead to the fixed cycle length. Numerical experiments are provided to test the performance of the proposed model using data generated by traffic simulation.

Index Terms—Connected Vehicles, Traffic Signal Optimization, Mixed Integer Nonlinear Program, Dynamic Programming, End Stage Cost, Branch and Bound.

I. INTRODUCTION

As a critical infrastructure that is crucial to the economy and the daily life of everyone, transportation also creates severe congestion and consumes tremendous energy. In the United States, e.g., the gasoline consumption by the transportation sector was about 143.37 billion gallons in 2016, a daily average of about 9.33 million barrels [1]. At the same time, traffic congestion on urban roads has caused extra fuel consumption as well as additional travel delays. The 2015 Urban Mobility Scorecard [2] estimated that U.S. highway congestion costs $160 billion a year, and an average American commuter loses 42 hours per year due to traffic congestion. Therefore, it is imperative to reduce traffic delay and improve transportation energy efficiency in urban areas.

Fuel consumption and traffic delay in urban areas can be reduced by optimizing traffic signal control strategies. Traditional traffic signal control problems have been extensively investigated, with a variety of methods such as fixed-time control, actuated control, and adaptive control [3]. The adaptive signal control methods (aka the most advanced traffic signal control methods so far), e.g., include swarm algorithm [4], platoon-based algorithms [5], rolling horizon approaches [6], and oversaturation algorithm [7], among others. A comprehensive discussion of the signal control algorithms can be found in [8].

Urban traffic signal control can be further enhanced by the connected vehicle (CV) technology [9]. CV enables vehicle to vehicle (V2V) and vehicle to infrastructure (V2I) communications through dedicated short-range communications (DSRC) and other means. With the V2V communications, location and speed information can be exchanged among nearby vehicles. With the V2I communications, vehicles can communicate with traffic signals, work zones, tollbooths, and other types of infrastructure to exchange information such as vehicle trajectories, traffic conditions, and signal timing. Such data/information exchange among vehicles and between vehicles and infrastructure has the potential to significantly improve traffic mobility and fuel consumption efficiency at signalized intersections.

There have been various traffic signal control studies under the CV environment. Goodall et al [10] developed the predictive microscopic simulation algorithm (PMSA) to control traffic signal. The strategy can minimize total delays, or the combination of delays, stops, and decelerations over a 15-second time period by considering instantaneous vehicle data. The study showed that at low or mid-level traffic volume, their proposed algorithm outperformed state-of-the-practice coordinated-actuated timing plan, while the performance got worse during saturated and oversaturated conditions. He et al. [11] developed the platoon-based arterial multi-modal signal control with online data (PAMSCOD) algorithm, with signal updated every 30 seconds. A mixed-integer nonlinear program (MINLP) was solved to determine future optimal signal plan. Simulation results in VISSIM showed that delays were significantly reduced under both non-saturated and oversaturated traffic conditions compared to traditional state-of-the-practice coordinated actuated signal control. Lee and Park [12] developed a cumulative travel-time responsive (CTR) real-time intersection control algorithm in the CV environment. They examined the different penetration rates of CV and level of congestion. Kalman filtering technique was utilized to estimate the cumulative of travel time under various penetration rates. Among various types of methods, dynamic programming (DP) is one of the most commonly used technique to solve the discretized signal control problems. It was first applied in Sen and Head [13] to optimize traffic signal timing. The idea was later applied in Chen et al. [14] and Feng et al.[15]. In particular, Feng et al. [15] proposed a bi-level formulation for optimizing
signal timing of a single intersection: the upper level is to optimize for the barrier lengths and the lower level is to optimize for the phase times. However, most of these studies assumed varying cycle lengths, which may not be readily applied to multiple intersections if signal coordination is needed. There are a few exceptions. He et al. [11] and Zamanipour et al. [16] incorporated fixed cycle length (thus coordination) as a virtual priority request, which introduces additional constraints in MILP. Beak et al. [17] extended Feng et al. [15] to impose the fixed cycle length, albeit with a bi-level formulation. First, they imposed extra constraints to the upper level to ensure a fixed cycle length. The revised intersection-level model (with a fixed cycle length) is then integrated into a corridor-level model for signal coordination. In this paper, we develop a two-step method to consider fixed cycle length at the intersection level. This avoids the use of the bi-level structure in Feng et al. [15] and Beak et al. [17], which is computationally more efficient.

Fuel consumption is a major consideration that researchers have been trying to evaluate when they develop traffic signal strategies. Zhao et al. [18] proposed a signal timing optimization strategy to minimize the combined total energy consumption and traffic delay, considering the fuel consumption of individual vehicles. Vehicles’ trajectories were predicted second by second using the Nagel-Schreckenberg model. An iterative grid search algorithm was used to solve the optimized signal timing. The method however ignored left-turn traffic and cannot be applied to real-world intersections. More studies that considered fuel prediction models can be found in [19]-[22]. Here we do not assume vehicles are autonomously driving, and therefore intersection control strategies, such as those based on reservation [23]-[26] and other methods [27]-[29], will not be discussed.

This study aims to optimize signal timing in the CV environment considering individual vehicle’s trajectories. As discussed above, most existing signal optimization methods work for single intersections and assume variable cycle lengths, with a few exceptions [16][17][30]-[32]. Coordination occurs when two or more traffic signals are working together so that moving vehicles could go through the intersections with minimum delays and the least number of stops. The cycle lengths of coordinated traffic signals are often fixed as a constant. In this study, we focus on a single intersection and assume a fixed cycle length such that the proposed method can be readily extended in the future to coordinate multiple signals in a traffic corridor or network. First, we formulate the CV-based signal control problem as a MINLP, by considering information of individual vehicle’s trajectories (i.e., second-by-second vehicle locations and speeds) and their realistic driving/car-following behavior, captured by the Intelligent Driving Model (IDM). The objective function is to minimize the weighted sum of total system fuel consumption and travel time. Due to the large dimension of the problem and the complexity of the nonlinear car-following model, solving the nonlinear program directly can be challenging. Secondly, we reformulate the problem as a DP model by dividing the timing decisions into stages (one stage for a signal phase) and approximating the total fuel consumption and travel time of a stage as functions of the state and decision variables. We note that imposing the fixed cycle length constraint will invalidate the DP formulation. We then apply a two-step method to resolve this issue. The first step is to add an end-stage cost to the DP formulation. The cost measures how much the DP solution violates the fixed cycle length constraint. This step forces the DP to produce a solution with a cycle length that is close enough to the given fixed cycle length. The second step is a branch and bound method to further refine the DP solution to obtain a solution of the original problem, with the exact fixed cycle length. Numerical experiments are provided in the paper to test the performance of the proposed model using data generated by traffic simulation.

The main contributions of the paper include:

1. This study accounts for the individual vehicle’s trajectories (i.e., second-by-second vehicle locations and speeds) into signal timing optimization, which can be provided by CV technologies.
2. A signal optimization strategy is developed for a single intersection with fixed cycle length, which can be easily extended for signal coordination. It is formulated as a mixed-integer nonlinear program (MINLP). The MINLP model may have a large dimension and is hard to solve.
3. A DP reformulation is proposed via certain approximation schemes. To ensure the fixed cycle length, a two-step method is developed: adding the end-stage cost and a branch and bound algorithm.

II. FORMULATING SIGNAL CONTROL AS A MIXED-INTEGER NONLINEAR PROGRAM

In this section, the signal control problem with the fixed cycle length constraint is formulated as a MINLP. Here we adopt the dual-ring method for signal design as it can properly balance safety and efficiency of traffic signal control [3][15]. This is important since the primary objective of traffic signal control is to ensure safety, i.e., to minimize conflicts [3], while mobility is also important as long as safety is ensured. The signal configuration in a dual-ring diagram is shown in Figure 1. Without loss of generality, we assume the eastbound/westbound (EB/WB) through movements (2 and 6 in Figure 1) are the major movements and thus cannot be skipped (i.e., for coordination purposes). Other phases may be skipped by setting the corresponding phase durations as zero. We also assume a cycle always starts with movements 2 and 6. Such a signal timing plan can be considered as 6 groups with a sequence of 8 phases in Figure 2. Noted that phase 2 and 3 in group 2 cannot be realized at the same time, indicating that at least one of the two phases need to be skipped. Same situation happens for phase 6 and 7 in group 5. In this paper, the continuous time is discretized into 1s intervals.

Based on the literature [33], the maximum/minimum green parameters are defined for a movement. In this paper, we define maximum/minimum green based on phases, which is consistent to [33] and Sen and Head [13] when phase overlaps are not allowed, i.e., only phase 1, 4, 5, and 8 exist, as shown in Figure 2. A phase overlap refers to a pair of phases that contain one common movement, e.g., phase 1 contains movement 2 and 6, and phase 2 contains movement 2 and 5, occurring in sequence would allow movement 2 to “overlap”. If phase overlaps exist, minimum greens for the overlap phases (phase 2, 3, 6, and 7)
are set to be zero in order to allow them to be skipped, the minimum greens for phases 1, 4, 5, 8 remain the same. Maximum greens for the overlap phases (2, 3, 6, and 7) are set to be a small value, e.g., 10s, while maximum green for the non-overlap phases need to be deducted by a corresponding value, e.g., the maximum green for phase 1 needs to subtract 10s. In the numerical test, we set the maximum green for non-overlap phases (1, 4, 5, and 8) as the cycle length, minimum green for non-overlap phases as 5s. For overlap phases (2, 3, 6, 7), the maximum green is half of the cycle length and minimum green is zero. This is designed to allow DP to search for sufficiently large state spaces to find an optimal solution, with a reasonable computational effort. The numerical tests suggest that selecting different values as maximum/minimum greens does not have significant influence on the total cost.

**Figure 1** Traffic signal configuration [35]

**Figure 2** Traffic signal configuration

### Parameters
- **C** Cycle length (s).
- **$m_F$** Monetary value of fuel ($/gal$), e.g., $3/gal$.
- **$m_{TT}$** Monetary value of travel time ($$/s$$), e.g., $12/h$ ($0.005/s$).
- **$e$** Idle fuel consumption rate (gal/h).
- **$l_n$** Length of vehicle $n$ (m).
- **$\delta$** Acceleration exponent in IDM. It usually set at 4.
- **$H$** Desired time headway (s), e.g., 1.5s.
- **$a_{max}$** Maximum acceleration rate (m/s²), e.g., 1 m/s².
- **$b_{max}$** Maximum deceleration rate (m/s²), e.g., 3 m/s².
- **$s_{0,n}$** Gap between vehicles in complete standstill traffic jams (m), e.g., 2m.
- **$v_0$** Vehicle desired speed (m/s).
- **$g_k$** Minimum effective green time of phase $k$ (s).
- **$g_k^{max}$** Maximum effective green time of phase $k$ (s).
- **$k$** Signal phase, $k=1,2...,8$.
- **$d_{n,t}$**, **$d_{signal,n}$** The location of the nearest front signal of vehicle $n$ (m).

### Variables
- **$FC_{s,n,t}$** Fuel consumption for vehicle $n$ at time $t$ (gal/s).
- **$TT_{n,t}$** Travel time of vehicle $n$ at time $t$ (s).
- **$FC_{i,n}$** Fuel consumption for vehicle $n$ at the idle status (gal/s).

The objective of the CV-based signal optimization problem can be formulated as minimizing the weighted sum of total fuel consumption and travel time [18][34] of all vehicles approaching the intersection:

\[
\min F = \sum_{n=1}^{N} \sum_{t=1}^{T} (m_F FC_{s,n,t} + m_{TT} TT_{n,t})
\]  

$FC_{s,n,t}$ and $TT_{n,t}$ are the fuel consumption and travel time for vehicle $n$ at time $t$. The corresponding parameters $m_F$ and $m_{TT}$ are the “value of fuel” and “value of time” respectively. Eq. (1) indicates that the objective function here considers the travel time and fuel consumption for all vehicles in the network for total time span $T$. $N$ is the total number of vehicles. Eq. (2) calculates the fuel consumption of vehicle $n$ at time $t$, which is determined by the vehicle status. If vehicle $n$ is idling at time $t$, the indicator variable for idle status, $I_{n,t}$, takes one and the fuel consumption model $FC_{s,n,t}$ is applied, as shown in Eq. (4). Otherwise, $I_{n,t}$ takes zero and Eq. (5) will be applied, which calculates the fuel consumption of vehicle $n$ at the moving status. The threshold value for idle status is 2.2 m/s (5mph), which is provided in Zhao et al [18]. Eq. (3b) reformulate (3a) using the “big $M$” method [36] by establishing a relationship between speed $v_{n,t}$ and idle status indicator $I_{n,t}$. $M$ here is a large number and $M=1,000$ is used in this paper. The model could be used to calculate fuel consumption for different vehicle types, including sedan, SUV, bus, EV, and Hybrid Electric (HEV). Zhao et al [18] provided the calibrated parameters in Eq. (4) and (5) for different vehicle types, as shown in Table 1.

\[
FC_{s,n,t} = FC_{s,n,t} \ast v_{n,t} \ast (1 - I_{n,t}) + FC_{i,n} \ast I_{n,t}
\]

\[
I_{n,t} = \begin{cases} 
1, & \text{if } v_{n,t} < 2.2, \text{ idling at time } t \\
0, & \text{if } v_{n,t} \geq 2.2, \text{ moving at time } t 
\end{cases}
\]

\[
\left\{ \begin{array}{l}
\frac{v_{n,t} - 2.2}{1 - I_{n,t}}M \\
\frac{v_{n,t} - 2.2 + I_{n,t}M}{M} \geq 0
\end{array} \right.
\]

\[
FC_{t,n} = e
\]

\[
FC_{s,n,t} = a \frac{v_{n,t}}{b + cv_{n,t} + dv_{n,t}^2}
\]
As aforementioned, this study assumes the cycle length be fixed for the whole time span T(e.g., a few hours). The effective green time for each phase k of cycle i, \(g_k^i\), must sum up to the (fixed) cycle length C, as shown in Eq. (6). We assume there is no transition time between phases. There are eight phases in Figure 1 so \(K = 8\). Eq. (7) indicates the bounds of the green time \(g_k^i\). For phases that can be skipped, \(g_k^\text{min} = 0\). Eq. (8-9) indicate phase 2 and 3 (and phase 6 and phase 7) cannot be realized for the same cycle i. At least one of the two variables, e.g., \(g_2^i (g_6^i)\) and \(g_3^i (g_7^i)\), needs to be zero.

\[
\sum_{k=1}^{K} g_k^i = C \quad \forall i \in 1, 2, ... I \quad (6)
\]

\[
g_k^\text{min} \leq g_k^i \leq g_k^\text{max} \quad (7)
\]

\[
g_2^i + g_3^i = 0 \quad (8)
\]

\[
g_6^i + g_7^i = 0 \quad (9)
\]

Each intersection contains multiple movements with two movements served by a phase k. Variable \(S_{k,t}\) denotes the signal status at time \(t\) for phase \(k\), as shown in Eq. (10a). Noted that when \(k = 0\), \(g_k^i = 0\). It takes one if the signal status at the current time stamp is red and zero if it is green. The variable \(k\) is the current phase index at time \(t\) (i.e., 1, 2... 8). It represents the phase in Figure 1 that is currently given the green light. Eq. (10b) reformulates (10a) using two binary variables \(y_{1,t}\) and \(y_{2,t}\) based on the big M concept.

\[
S_{k,t} = \begin{cases} 
0, & \text{if } \sum_{k=0}^{K-1} g_k^i \leq (t \mod C) < \sum_{k=0}^{K-1} g_k^i + M \\
1, & \text{otherwise}
\end{cases} \quad (10a)
\]

\[
\sum_{k=0}^{K-1} g_k^i \leq (t \mod C) + y_{1,t}M > 0 \\
\sum_{k=0}^{K-1} g_k^i - (t \mod C) \leq (1 - y_{1,t})M \\
(t \mod C) - \sum_{k=0}^{K-1} g_k^i < (1 - y_{1,t})M \\
(t \mod C) - \sum_{k=0}^{K-1} g_k^i < (1 - y_{1,t})M
\]

\[
S_{k,t} = y_{1,t} + y_{2,t} \quad (10b)
\]

Eq. (11) use indicator variables \(y_{1,t}\) and \(y_{2,t}\) together to identify whether vehicle \(n\) is within the “range” of the incoming approach defined by \(d_{\text{in}}^n\) and \(d_{\text{in}}^\text{min}\). Both of them will be zero if vehicle \(n\) is within the range; otherwise, one of the will be 1. Furthermore, the signal status \(Z_{n,t}\) at time \(t\) for vehicle \(n\) could be determined as long as the incoming approach of vehicle \(n\) is identified. There are usually two phases that may serve vehicle \(n\), e.g., both phase 1 and phase 2 could serve vehicle \(n\) if it comes from movement 2. In this case, we first define \(Z_{n,t}\), the “minimum” signal status of the two phases that serve vehicle \(n\); see equation (12a) below. This is to ensure that as long as one of the two phases is green \(S_{k,t}\) is zero), vehicle \(n\) will see green \(Z_{n,t}\) is zero). Note that signal status \(Z_{n,t}\) and \(S_{k,t}\) are different.

Before the vehicle enters the intersection range, the trajectory of the vehicle should not be impacted by the signal. This is guaranteed by Eq. (12b) that \(Z_{n,t}\) is always zero (green) before the vehicle enters the intersection range so that neither the location nor the speed of the vehicle is impacted by \(Z_{n,t}\); see equations (14) and (15). Once the vehicle enters the range, \(Z_{n,t}\) is the same as \(Z_{n,t}\) as shown by equations (11) and (12b), which also equals to one of the \(S_{k,t}\)’s as shown in (12a).

\[
Z_{n,t} = \begin{cases}
\min\{S_{k,t}\} & \text{for all phase } k \text{ that may serve vehicle } n \quad (12a)
\end{cases}
\]

The CV-based signal timing strategies in this paper require information on real time vehicle trajectories. This study assumes the 100% penetration rate of CVs. Vehicle trajectories can be transmitted when a vehicle enters the boundary of an intersection. Furthermore, to optimize signal timing for the current and future cycles, future vehicle trajectories are needed. For this, the Intelligent Driver Model (IDM) [37] is applied to simulate vehicle trajectories. IDM is a car-following model that fits better with CV [38]. We assume that there is only one lane (and a dedicated left turn lane at the intersection) per incoming approach, so there is no lane changing behavior involved. It is necessary to account for the signal status in the prediction of traffic flow propagation when applying IDM. For this, we model the red signal as a “standing vehicle” with speed equal to zero. It would disappear if the signal turns green. Eq. (13a) indicates whether the front object of vehicle \(n\) is a real vehicle or a standing vehicle (traffic signal) by comparing the relative location of the front vehicle \(n - 1\), vehicle \(n\), and the nearest traffic signal in front of vehicle \(n\). The binary variable \(y_{n,t}\) takes one if the front “vehicle” is the traffic signal (could be red or green) at location \(d_{\text{signal}, n}\) with speed zero. If \(y_{n,t}\) is zero, the front vehicle \(n - 1\) is a real vehicle with location \(d_{n-1,t}\) and speed \(v_{n-1,t}\). This helps update the vehicle trajectories in IDM as shown later. Eq. (13b) reformulate (13a) using the big M method and two binary variables \(y_{n+1,t}\) and \(y_{n+2,t}\).

\[
y_{n,t} = \begin{cases} 
1, & \text{if } x_{n,t} < d_{\text{signal}, n} < x_{n-1,t} \\
0, & \text{otherwise}
\end{cases} \quad (13a)
\]

\[
\begin{align*}
&d_{\text{signal}, n} - d_{n-1,t} < y_{n,t} < d_{\text{signal}, n} - d_{n-1,t} + (1 - y_{n,t})M \\
&d_{n,t} - d_{\text{signal}, n} < (1 - y_{n,t})M
\end{align*} \quad (13b)
\]

Using IDM, Eq. (14-15) identify the vehicle location and speed of the preceding “vehicle” \(n - 1\), which could be a real vehicle or the nearest front signal.

\[
f_{n-1,t}^d = v_{n-1,t} * \left[ \frac{1}{2} y_{n,t} + x_{n,t} - \frac{1}{2} y_{n,t} + x_{n,t} - Z_{n,t} \right] + d_{\text{signal}, n,t} \quad (14)
\]

\[
f_{n-1,t}^v = v_{n-1,t} * \left[ \frac{1}{2} y_{n,t} + x_{n,t} - \frac{1}{2} y_{n,t} + x_{n,t} - Z_{n,t} \right] \quad (15)
\]
Eq. (16-19) shows how IDM estimates the acceleration rate for vehicle \( n \) at each time interval, given the location and speed of vehicle \( n-1 \).

\[ s_{n,t} = F_{n-1,t}^d - d_{n,t} - l_{n-1} \]  \tag{16}

\[ \Delta v_{n,t} = f_{n-1,t}^v - v_{n,t} \]  \tag{17}

\[ a_{n,t} = a_{\text{max}} \left[ 1 - \left( \frac{v_{n,t}}{v_0} \right)^\delta - \left( \frac{s'(v_{n,t})\Delta v_{n,t}}{s_n^2} \right) \right] \]  \tag{18}

\[ s'(v_{n,t},d_{n,t}) = s_n + v_{n,t}H + \frac{2}{2a_{\text{max}}^2} \]  \tag{19}

Eq. (20-21) are applied to update the trajectories for vehicle \( n \) at the next time interval \( t+1 \). More details of Eq. (14-21) can be found in [37]. In this paper, the values of the parameters in IDM are chosen as \( a = 1 \text{m/s}^2, b = 3 \text{m/s}^2, s_o = 2 \text{m}, H = 1.5 \text{s}, \) and \( \delta = 4 \), according to [39].

\[ v_{n,t+1} = \max(0, v_{n,t} + a_{n,t}) \]  \tag{20}

\[ d_{n,t+1} = d_{n,t} + \frac{v_{n,t} + v_{n,t+1}}{2} \]  \tag{21}

Eq. (1) – (21) is a MINLP for the CV-based signal control problem. It clearly shows that when individual vehicle status is considered for signal control, e.g., under the CV environment, the problem can be formulated as a very complex MINLP. This is mainly due to the different status of vehicles and signal phases, as well as the various if-then-else types of conditions (e.g., equations (3), (8), (13), and others) inherent to this coupled signal-vehicle optimization problem. In addition, since the variables of the model include the location and speed of each vehicle at each time interval, the dimension of the problem can be quite large. Furthermore, the IDM-based car following model is also very complex. Thus solving the model directly is quite challenging, and more tractable and efficient methods are needed. We next present one of such methods based on DP.

III. DYNAMIC PROGRAMMING FORMULATION

DP provides a general framework to divide an optimization problem into multiple stages (under certain conditions), which could be solved sequentially one stage at a time. Here we divide the signal timing decisions into stages, one stage for a phase. We then approximate the total fuel consumption and travel time of a stage as functions of the state and decision variables of that stage only. The notation for DP is first summarized as follows; the unit of each variable/parameter is provided in parentheses.

\( x_p \) Decision variable, phase duration of stage \( p \) (s).

\( s_p \) State variable, total time from beginning of the cycle to the end of stage \( p \) (s).

\( X_p(s_p) \) The set of feasible control variable given stage variable \( s_p \) at stage \( p \) (s).

\( V_p(s_p) \) Value function, the cumulative value of objective function from stage 1 up to stage \( p \) ($).

\( x_{p}^{\min} \) Minimum value of the decision variable at stage \( p \) (s).

\( x_{p}^{\max} \) Maximum value of the decision variable at stage \( p \) (s).

\( f_p(s_p,x_p) \) Total cost at stage \( p \), given state variable \( s_p \), and decision variable \( x_p \) ($).

\( N_p \) Total number of vehicles in phase \( p \) (veh).

\( FC_{n,t}(s_p,x_p) \) Fuel consumption of the vehicle \( n \) at time \( t \) given stage variable \( s_p \) and decision variable \( x_p \) (gal/s).

\( TT_{n,t}(s_p,x_p) \) Travel time of vehicle \( n \) at time \( t \) given stage variable \( s_p \) and decision variable \( x_p \) (s).

\( FC_{n,t+1}(s_p,x_p) \) Fuel consumption of vehicle \( n \) at moving status at time \( t \) given stage variable \( s_p \) and decision variable \( x_p \) (gal/meter).

\( FC_{n,t}(s_p,x_p) \) Fuel consumption of vehicle \( n \) at idle status at time \( t \) given stage variable \( s_p \) and decision variable \( x_p \) (gal/s).

\( A_p(s_{p-1},s_p) \) The number of arriving vehicles in the time interval \([s_{p-1},s_p])\).

\( M_p(x_p) \) The maximum number of vehicle that can be discharged during phase duration \( x_p \).

\( v_{n,t}(s_p,x_p) \) Approximated speed of vehicle \( n \) at time \( t \) given the stage variable \( s_p \) and decision variable \( x_p \) (m/s).

\( t_a \) Arrival time at the predefined intersection boundary (distance \( L \) upstream of intersection) (s).

\( t_d \) Time when vehicle joins the queue (started to slow down) (s).

\( t_0 \) Time when the vehicle fully stops (s).

\( t_{dc} \) Time when vehicle starts to be discharged (s).

\( t_1 \) Time when the vehicle achieves the free flow speed \( v_f(s) \) (s).

\( l_d \) Distance upstream of end of queue or the stop line (if no queue) (m), e.g., 100m.

\( V_p(s_p) \) Value function at phase \( p \) given state variable \( s_p \).

\( \sigma \) Tolerance of the fixed cycle length (s), e.g., 5s.

As shown in Figure 1 and Figure 2, there are eight stages in total. The state variable \( s_p \) is defined as the total number of time intervals from the beginning of the cycle to the end of stage \( p \), while the decision variable \( x_p \) is the phase duration. Eq. (22-23) illustrate the relationship between state variable and decision variable. Parameter \( r \) denotes the effective clearance interval; see Sen and Head [13] for more details.

\[ s_p = s_{p-1} + h(x_p) \]  \tag{22}

\[ h(x_p) = \begin{cases} x_p + r, & \text{other wise} \end{cases} \]  \tag{23}

Given the state variable \( s_p \), the feasible set of decision variables could be determined based on Eq. (24):

\[ X_p(s_p) = \begin{cases} 0, & \text{if } s_p < x_p^{\min} \\ \{0,x_p^{\min}, x_p^{\min} + 1, ..., (\min(x_p^{\max}, s_p)) \}, & \text{otherwise} \end{cases} \]  \tag{24}

To formulate the DP, we first assign the initial value function \( V_0 = 0 \). The DP starts from stage (phase) \( p = 1 \), and proceed recursively to \( p = 2, 3 \ldots 8 \). At each stage, the method calculates the optimal decision variable \( x_p^*(s_p) \) by minimizing the value function for each possible value of the stage variable \( s_p \) in the forward recursion. Note that after stage 1, phase 2 or phase 3 may be chosen for the optimal decision variable, but not both; the same rule applies to phase 6 and
phase 7. After the decision variables are estimated for all stages, the optimal decision of each stage can be retrieved in the backward recursion. The DP calculation process can be equivalently represented by an acyclic graph, as shown in Figure 3.

Figure 3 Acyclic graph of DP calculation process

However, in order to reformulate the signal control problem (1) – (21) as a DP, a critical condition is that the objective function in (1) can be expressed as the summation of the objective function of each stage. Furthermore, the stage-specific objective function (i.e., the sum of the vehicle fuel consumption and travel time of all vehicles in the stage) can be expressed as a function of the state and decision variables of that stage only [13]. This however is not true in general for most of the objectives we consider here, i.e., travel time or fuel consumption. It is especially so when we consider the data/information of individual vehicles (such as trajectories, speeds, delays, etc.). In the next subsection, we approximate the objective function of each stage so that it can be expressed as a function of the state and decision variables of the stage.

A. Objective Function Approximation

Eq. (25a) expresses the total fuel consumption and travel time of all the vehicles for phase $p$ (i.e., it is from time $s_{p-1}$ to $s_p$), where $N_p$ is the total number of vehicles in phase $p$. In this paper, travel time of a vehicle is estimated by the summation of free flow travel time of the vehicle and the delay it encountered. As shown previously [13], the total delay (and thus travel time) of a stage in (25b) can be approximated as a function of the state and decision variables. We show in this subsection how the fuel consumption can be approximated as a function of the state and decision variables. As shown in (5) and rewritten in (26), fuel consumption is a function of vehicle speeds. Thus we aim to approximate the vehicle speed as a function of the state and decision variables.

$$
\text{Min } \sum_{n=1}^{N_p} \sum_{s_{p-1}}^{s_p} f_p(s_p, x_p) = \sum_{n=1}^{N_p} \sum_{s_{p-1}}^{s_p} f_p(s_p, x_p) = m_p FC_{n,t}(s_p, x_p) + m_{TT,n,t}(s_p, x_p)
$$

The speed of a vehicle is estimated based on the queue discharging process. Let $A_p\left(s_{p-1}, s_p\right)$ denote the number of arriving vehicles for phase $p$, i.e., during the time interval $[s_{p-1}, s_p]$. $M_p\left(x_p\right)$ denotes the maximum number of vehicles that can be discharged during the time interval of green time $x_p$.

As shown in [13], $A_p$ and $M_p$ can be expressed as functions of the state and decision variables of stage $p$ only. We next show how the speed of a vehicle can be approximated as those variables. There are four possible cases for the speed of a vehicle arriving in stage $p$, as shown in Figure 4.

1. Vehicle arrives during green signal at stage $p$ and can pass freely through intersection.
   $$v = v_0$$

2. Vehicle arrives during green but a queue already exists. The queue includes vehicles that were not discharged in the previous stage $p-1$ and the newly arrived vehicles at stage $p$ before the current vehicle. Denote $t_d$ the time when the current vehicle arrives at the predefined intersection boundary (distance $l$ upstream of intersection), $t_d$ the time when the vehicle joins the queue (started to slow down), $t_0$ the time when the vehicle fully stops, $t_{ac}$ the time when the vehicle starts to be discharged, and $t_1$ the time when the vehicle achieves the free flow speed $v_0$ again after being discharged. We assume the vehicle starts to decelerate with a constant rate at time $t_d$ if queue exists. We can then approximate the average speed between $t_d$ and $t_0$ as $\frac{v_0}{2}$. The same assumption applies to the acceleration process from $t_{ac}$ to $t_1$. When the vehicle starts to decelerate at $t_d$, the distance between the vehicle and the stop line is denoted as $l_d$.

   If the vehicle could pass the intersection within the current phase $p$ (trajectory (1) in Figure 4(b)), speed could be approximated using Eq. (28), otherwise (trajectory (2)), vehicle has to wait until the next phase. In this case, we only apply first three conditions in Eq. (28) since $t_1 \geq s_p$. We could consider (2) as a special case of (1).

   $$v_{n,t}(s_p, x_p) = \begin{cases} 
   v_0, & \text{if } t_d \leq t \leq t_{ac} \\
   \frac{v_0}{2}, & \text{if } t_d < t \leq t_0 \\
   0, & \text{if } t_0 < t \leq t_{ac} \\
   \frac{v_0}{2}, & \text{if } t_{ac} < t \leq t_1 \\
   v_0, & \text{if } t_1 < t \leq s_p 
   \end{cases}$$

3. Vehicle arrives during red signal at stage $p$. Figure 4(c) and (d) are differentiated by whether queue exists at the end of stage $p$. If there is no queue, vehicle $n$ will start to leave at $t_1\left(t_1 = s_p\right)$, otherwise it will have to wait for the queue to dissipate ($t_1 \geq s_p$). Here we only care about the vehicle status from time $s_{p-1}$ to $s_p$. The speed of vehicle $n$ arriving during red can be summarized as:

   $$v_{n,t}(s_p, x_p) = \begin{cases} 
   v_0, & \text{if } t_d \leq t \leq t_{ac} \\
   \frac{v_0}{2}, & \text{if } t_{ac} < t \leq t_0 \\
   0, & \text{if } t_0 < t \leq s_p 
   \end{cases}$$

After approximating the vehicle speed $v_{n,t}(s_p, x_p)$ in the objective function for stage $p$ based on the above four cases, the objective function can be approximated as a function of $s_p$ and $x_p$ only, as shown below:

$$v_p(s_p) = \min\{f_p(s_p, x_p) + V_{p-1}(s_{p-1}) | x_p \in X_p(s_p), p \in P\}$$

This means that Eq. (25), and the objective function of the MINLP in Section II as well, can be expressed as a function of $s_p$ and $x_p$ (for all $p$) only. Thus the problem can be reformulated as a DP. Detailed proof is straightforward and
omitted here. Furthermore, after obtaining the optimal decision variable $x_p$ at stage $p$, IDM is applied to update the trajectories of all the vehicles from $s_{p-1}$ to $s_p = s_{p-1} + x_p$.

$$\text{Figure 4 Vehicle average speed Approximation}$$

(a) Vehicle arrives during green signal (no queue)
(b) Vehicle arrives during green signal and queue exists
(c) Vehicle arrives during red signal (no queue)
(d) Vehicle arrives during red signal and queue exists

This constraint however will invalidate the DP model as decisions cannot be made stage by stage when Eq. (34) is considered. In other words, DP cannot guarantee that the optimal phase durations sum up to a fixed cycle length. A branch and bound method is applied here to resolve the issue, which was used in the past to solve the Resources Constrained Shortest Path (RCS) problem [40]. The method creates a tree by selecting one variable each time from an initial solution. Here the initial solution is produced by solving the DP formulation with the end-stage cost, as discussed above. The maximum level of the tree is the number of stages because the variable in a given set can only be used for branching once. Moreover, all branches at a given level of the tree have to be computed and analyzed before advancing to the next level. The numerical example of the branch and bound method will be provided in the next section. The can be summarized as below:

1. **DP with the end-stage cost is solved first to produce an initial solution.**
2. **Define the error gain for stage/phase $p$:**
   $$E_G(p, s_p, x_p) = f_p(s_p, x_p)$$
   (35)

   The error gain is the objective function in DP (total fuel consumption and travel time) for all vehicles traveling in stage $p$ with stage variable $s_p$ and decision variable $x_p$.
3. **If the estimated cycle length happens to be exactly the predefined cycle length $C$, the DP solution is an optimal solution. The algorithm stops.**
4. **If the cycle length from the DP solution is larger than the predefined fixed cycle length (decision variable should be decreased), the selected phase for branching at each level is the one that has the minimum error gain.** There could be multiple number of branching depending on the difference between the predefined cycle length and the cycle length from the DP solution. If the produced cycle length is less than the fixed cycle length (decision variable should be increased), the selected phase for branching would be the maximum error gain.
5. **The algorithm stops when the results of all branching (i.e., leaves) are feasible. The optimal solution is selected from the feasible solutions that satisfies the cycle length constraint while producing the minimum objective value.**

**IV. NUMERICAL EXPERIMENTS**

We test the proposed CV-based signal timing optimization model and the DP method using data generated in VISSIM, which include the arrival time, arrival link (movement) and initial speed of each vehicle. The testing network contains a single intersection with a boundary of 300 meters upstream and downstream of the intersection to mimic the communication range of V2I. After comparing the approximated vehicle speed using our approximation method and the speed generated from VISSIM, we evaluate the proposed signal timing optimization method in three steps. First, we estimate the optimal cycle length using SYNCHRO, a widely used traffic signal design and optimization software tool, for different traffic demand cases. Next, for a given case, we apply different methods (see Table 2 and Table 3) to optimize the signal timing plans (phases and green splits), one for each method. We then evaluate the performance of each signal timing plan (i.e., each method) by applying the plan to the intersection and using IDM to generate
vehicle trajectories, based on which to calculate the total cost. We also illustrate the procedure of the branch and bound algorithm and test the impact of the tolerance in the end-stage cost Eq. (32) on the performance of the proposed model. The values of all parameters used in this study can be found in the notation lists.

A. Speed approximation

In the proposed method, we approximate the vehicle speed as a function of the state variables and decision variables, as shown in Section III. Here we compare the vehicle speed between our method (Eq. (27)-(29)) and IDM simulation for the major streets in Figure 5. The suggested value of parameters in IDM are indicated in the notation lists. The platoon contains four vehicles. It is observed that the approximate speeds have the similar trend but not as smooth as the speed profile in IDM.

![Figure 5 Speed Comparisons](image)

(a) IDM speed  (b) Approximated speed

In order to validate the speed approximation method is acceptable in fuel consumption estimation, the absolute differences of the fuel consumption with and without the speed approximation are estimated. It is observed that under all six cases, the mean absolute errors of fuel consumption using approximate speed are all less than 16%, which are considered acceptable in fuel consumption estimation.

B. Signal Timing Optimization

Different combinations of traffic demands and vehicle types are tested in order to evaluate the proposed signal optimization method. Vehicle arrivals at the boundaries of the intersection are generated from VISSIM. For each direction, 80% of the vehicles will go straight and the others will turn left. Six cases are tested to reveal the influence of traffic demands and vehicle types on the model performance. In Case I – III, vehicle demand is set to be 250 vph, 500 vph, and 800 vph, respectively. All vehicles are sedans (Type 5 in Table 1). In Case IV - VI, traffic demands are identical as in Case I - III, but the vehicle types are assigned differently. In the N-S directions, vehicles are assigned as Electric Vehicles (EVs, i.e., Vehicle type 1), while in the W-E directions, vehicles are assigned as buses (Vehicle type 7). Notice that generating all buses from one direction and all EVs from another is not very realistic. This is done here mainly to show more clearly the influences of the vehicle types on fuel consumption and signal timing.

We first test all models for 10 cycles. There are vehicles randomly generated to enter the network during the first 8 cycles, while during the last 2 cycles, there is no traffic demand. This guarantees the network is cleared by the end of the simulation. We test three signal control methods, as shown in Table 2 and Table 3. The first method is the actuated signal plan every cycle. The results from DP without fixed cycle length constraints are also shown in the tables.

Table 2 shows the costs from different models under various demand levels. NOMAD produces different solutions for different signal plan updating intervals and the starting points (from SYNCHRO in first four rows and from DP in the last row). We use the DP solution or the Synchro results as the initial points for NOMAD to make sure it can start with some reasonably good initial points. The best solutions from NOMAD are highlighted for each Case.

Table 2 shows that if the signal plan is updated every cycle or every two cycles, the network performance are not improved since the total costs keep the same as its initial evaluation from the initial guess (SYNCHRO plan). However, a better solution may be obtained as the updating frequency decreases to every 5 or 10 cycles. This indicates that NOMAD has difficulties finding optimal solutions when the number of variables is relatively large. Table 2 also shows that different starting points in NOMAD can affect the optimal solutions. When compared with different models, DP outperforms SYNCHRO and achieves the same solution as NOMAD under low and medium demands in case I and case II. In case III, as the demand increases, DP results are still better than SYNCHRO but slightly worse than NOMAD. Table 3 shows the costs from different models under various demand levels and vehicle types. By comparing Case IV to Case I, we can see that the performance improvements of DP and NOMAD over SYNCHRO are more dramatic, which is also shown in Figure 6. Case IV incorporates buses in W-E direction, which produces much more fuel consumption than sedan in Case I. The cost generated from DP and NOMAD were both lower than SYNCHRO because SYNCHRO does not consider the
influence of vehicle types on the fuel consumption when optimizing the signal timing. As the demand increases while still maintaining different vehicle types on NS and WE directions, the performance of DP is still better than SYNCHRO but slightly worse than NOMAD in medium (500 vph) and high (800 vph) demand levels. Since NOMAD uses the optimal phase plan from SYNCHRO and DP as the starting point, it makes sense that the best NOMAD results are always better than or equal to SYNCHRO and DP for all cases. It is also observed that the results from DP without the fixed C constraint always have the lower objectives compared with the one with the fixed C constraint. This indicates that enforcing the fixed cycle length, i.e., for signal coordination purposes (and thus to the benefit of the entire system), may likely worsen the performance of individual intersections.

Table 2 Total Cost under Various Demand Levels

<table>
<thead>
<tr>
<th>Model</th>
<th>Case I</th>
<th>Case II</th>
<th>Case III</th>
</tr>
</thead>
<tbody>
<tr>
<td>SYNCHRO</td>
<td>49.64</td>
<td>162.39</td>
<td>368.99</td>
</tr>
<tr>
<td>NOMAD update every 5 cycles</td>
<td>49.64</td>
<td>162.39</td>
<td>368.99</td>
</tr>
<tr>
<td>NOMAD update every 10 cycle</td>
<td>48.12</td>
<td>143.52</td>
<td>358.05</td>
</tr>
<tr>
<td>NOMAD update every 10 cycles &amp; initial points from DP solution</td>
<td>47.74</td>
<td>140.65</td>
<td>359.96</td>
</tr>
<tr>
<td>DP with fixed C constraint</td>
<td>47.74</td>
<td>140.65</td>
<td>360.73</td>
</tr>
<tr>
<td>(DP without fixed C constraint)</td>
<td>46.97</td>
<td>138.55</td>
<td>354.74</td>
</tr>
</tbody>
</table>

Figure 6 shows the performance improvements of NOMAD, DP with fixed cycle length, and DP without fixed cycle length over SYNCHRO. As shown in the dashed line, the model improvements of low and high demand levels are not as significant as the middle demand levels. This may be because under unsaturated but relatively heavy traffic conditions, there are more opportunities to optimize the splits and reduce the total cost of fuel consumption and travel time. Such opportunities tend to diminish when traffic is very light (all methods can work well) or very heavy (no method can work well). Furthermore, the performance improvements are more obvious if considering different vehicle types, as shown in Case IV – VI in Figure 6.

Figure 7 shows the cost of fuel consumption and travel time separately for the four methods and six cases. For all cases, the cost of travel time is much larger than the cost of fuel consumption. Comparing cases IV, VI to cases I, II, III, it is observed that the influence of vehicle types is more significant on the cost of fuel consumption than travel time. Considering the same level of travel demand (e.g., case I and IV), the cost of fuel consumption is larger for the cases considering different vehicle types while the costs of travel time are similar.

The computation times for NOMAD and DP depend on the several factors, e.g., the update intervals of signal plans, whether considering vehicle types and traffic demands in the network. For NOMAD, the initial points and number of variables also influence the optimization time. Under high traffic demand, NOMAD can take up to 20 minutes to find a solution if signal timing is updated every 10 cycles; in some cases, NOMAD may fail to find any solution. For DP, the optimization time ranges from 10-20s per cycle (in the cycle-by-cycle signal optimization), depending on the traffic demand level. Currently the methods are implemented in Matlab and no code optimization is performed. When signal timing optimization is conducted in real time, the methods will be most likely programed in a more efficient language/platform (such as C or machine language) and the codes should be optimized to improve the performance. This can often significantly improve the computation time. Detailed investigation of this issue is beyond the scope of the paper, and may be pursued in future research.

In order to further verify the algorithm, the simulation period is extended from 10 cycles to 1 hour. For NOMAD, it cannot find a feasible solution even if the updating interval for signal timing is 10 cycles (larger updating intervals indicate fewer number of variables in NOMAD). Therefore we consider NOMAD fail to solve the signal optimization problem for the 1 hour simulation period and results are not shown here. In Table 4, the cost estimated from DP is less than SYNCHRO in every case. Case V has the largest improvement (10.02%) when considering various vehicle types under medium vehicle demands.
Table 4 Cost of Different models for 1 hour simulation

<table>
<thead>
<tr>
<th>Model</th>
<th>Case I</th>
<th>Case II</th>
<th>Case III</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>250 vph, sedan</td>
<td>500 vph, sedan</td>
<td>800 vph, sedan</td>
</tr>
<tr>
<td>1. SYNCHRO</td>
<td>277.72</td>
<td>685.39</td>
<td>1388.55</td>
</tr>
<tr>
<td>2. DP</td>
<td>Improvements (%)</td>
<td>266.39</td>
<td>4.08%</td>
</tr>
</tbody>
</table>

C. Branch and Bound Algorithm

This section illustrates how the branch and bound algorithm can be applied to the DP results (without the fixed cycle length constraint) to guarantee a solution with the fixed cycle length. Here we use the DP result of Case I in Table 2. Figure 8 shows the result from the DP by considering the end stage cost, which leads to a cycle length of 57s, 3 seconds lower than the predefined and fixed cycle length of 60s.

D. Tolerance parameter of branch and bound method

In Eq. (32), we define a tolerance \( \sigma \) to calculate the end stage cost to find the DP solution. Different values of \( \sigma \) may produce different initial solutions and influence the number of evaluations in the Branch and Bound algorithm. Here an evaluation means that for a given tentative signal timing plan, we need to estimate the objective in (25) using IDM, which may be time consuming. As shown in Figure 9, one node in the graph corresponds to one signal timing plan that needs to be evaluated. We test the tolerance \( \sigma \) from 0 to 10 and test the performance of the algorithm for Case I. In Figure 10, as \( \sigma \) increases, the total cost decreases slightly and attains its minimum value at \( \sigma = 8 \), but the number of evaluations increases dramatically. A very large value of \( \sigma \) (e.g., 10) will need to evaluate more timing plans, but may not help much to minimize the objective. Similar trends can be found for other cases. Setting \( \sigma = 5 \), as we use in this paper, seems a good balance between solution quality and the computational effort of the algorithm.

Figure 9 Branch and Bound Tree

![Branch and Bound Tree](image)

Figure 10 Influence of \( \sigma \) on the total cost for Case I

![Influence of \( \sigma \) on the total cost](image)

V. CONCLUSION

This paper presents a signal timing optimization model for a single intersection with a fixed cycle length under the CV environment. The algorithm utilizes arrival information (speeds, locations, etc.) from CV as the input to optimize the green time by considering individual vehicles’ fuel consumption and travel time. The problem was first formulated as a mixed-integer nonlinear program (MINLP) by applying the...
IDM to predict vehicle trajectories. Such a formulation has a large dimension and a complex car-following model (the IDM). A DP formulation was then developed to approximate the MINLP. The overall problem was divided into stages (one stage for each signal phase). The objective is the summation of the objective of each stage. The objective function of a stage was approximated as a function of the state and decision variables of the stage only, by approximating the vehicle speeds and delays. We showed that imposing the fixed cycle length constraint would invalidate the DP formulation. We then applied a two-step method to address this issue. First, we added an end-stage cost to the DP formulation, defined by how much the DP solution violates the fixed cycle length constraint. This step forced the DP to produce a solution with a cycle length that is close to the given fixed cycle length. The second step was a branch and bound method to further refine the DP results to obtain a solution that produces the given cycle length exactly.

We evaluated the performance of the algorithm using data generated from traffic simulation. The results of the proposed DP model were compared with two other models. The first one is the traditional actuated signal timing plan generated by SYNCHRO. The second is to solve the MINLP formulation directly using the NOMAD solver in MATLAB. The results showed that the proposed DP method is always superior to SYNCHRO under all cases and can generate similar (slightly worse) solutions compared with NOMAD. However, NOMAD has difficulties finding optimal solutions when the number of variables is relatively large and the computational times of NOMAD are much larger than DP. This makes the proposed DP method more favorable when dealing with large scale problems.

Overall, the obtained solution by the proposed model ensures the (given) fixed cycle length, which is crucial for extending the proposed method to optimize and coordinate multiple traffic signals in a traffic corridor or network in the future. The objective for such corridor level optimization and coordination is to produce optimal offsets by minimizing the total fuel consumption and travel times of vehicles traveling along the coordinated movements [41]. For this, the proposed single-intersection optimization method, especially the DP reformulation and the two-step method, serves as a crucial component. The authors are investigating such a signal coordination problem and results may be reported in subsequent papers. Future research may also investigate how different penetration of CV-equipped vehicles will affect the performance of the proposed signal control method. This will require estimating the trajectories of vehicles that are not equipped with CV. When sample trajectory data from real world are available, certain stochastic models, e.g., Kalman filter based or Bayesian methods may be applied to estimate and predict the trajectories. For this, past work of the authors on estimating vehicle trajectories at signalized intersections may be helpful [42]. Recent study by Continental [43] can also be insightful for this. They applied the complete sensor set for an intersection to track vehicle trajectories, which can largely improve the penetration rate of CV. Furthermore, the proposed method needs to be tested using real world traffic signals and CV data. This will be pursued in future research.

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